

Scalable Fault Tolerant Protocol for Parallel Runtime Environments

Thara Angskun, Graham E. Fagg, George Bosilca,
Jelena Pjesivac-Grbovic, and Jack J. Dongarra

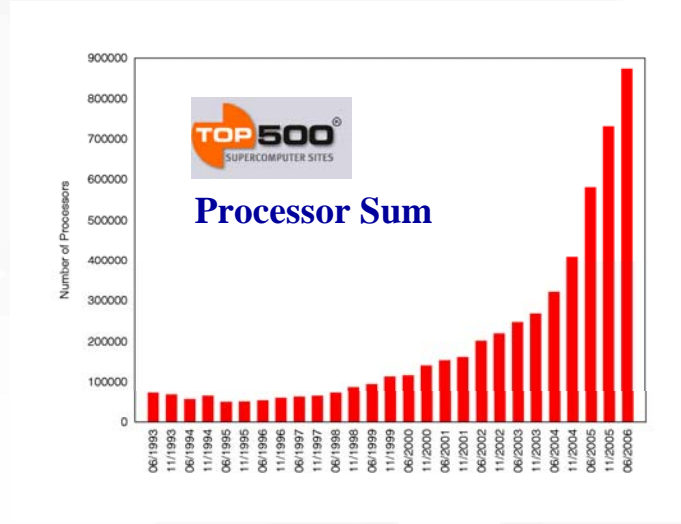
Euro PVM/MPI 2006
(09/19/06)



INNOVATIVE COMPUTING LABORATORY

COMPUTER SCIENCE DEPARTMENT
UNIVERSITY OF TENNESSEE

- Introduction
- Background
- Design
- Verification
- Results
- Conclusion



- » Increase number of processors
- » Decrease MTTF
 - » Dynamic Environment

- » Parallel Runtime Environment
 - » Extension of OS services for message passing library or application development
 - » **SCALABLE** and **FAULT-TOLERANT**

Introduction

Background

Design

Verification

Results

Conclusion

- » MPI runtime environments
 - » Start / terminate jobs
 - » Transfer signals (e.g. Ctrl-C)
 - » Redirect STDIN, collect stdout / stderr
 - » Collect exit status
 - » Monitoring job status
 - » (Optional) Interface with debugger, scheduler etc.

- » Communication Protocol
 - » Handle multiple types of message transmissions
 - » Broadcast, Multicast, Unicast
 - » **SCALABLE** and **FAULT-TOLERANT**

Introduction

Background

Design

Verification

Results

Conclusion

- » MPICH2 - MPD (Multi-Purpose Daemon)
 - » Ring or Tree topology

- » Open MPI – Open RTE
 - » Linear

- » LAM/MPI – LAM
 - » Linear

- » FT-MPI – HARNESS
 - » Linear

Introduction

Background

Design

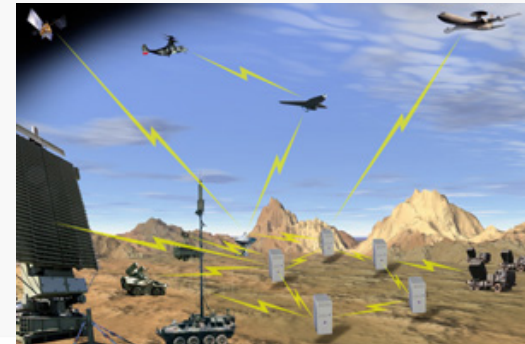
Verification

Results

Conclusion

- » Structured peer-to-peer networking
 - » Based on distributed hash tables
 - » CAN, Chord, Pastry, Tapestry
 - » Focus on resource discovery (Unicast)

- » Sensor or large scale ad-hoc networking
 - » Based on gossiping (epidemic algorithm)
 - » Focus on information aggregation.



Scalable and Fault-Tolerant Protocol

Introduction

Background

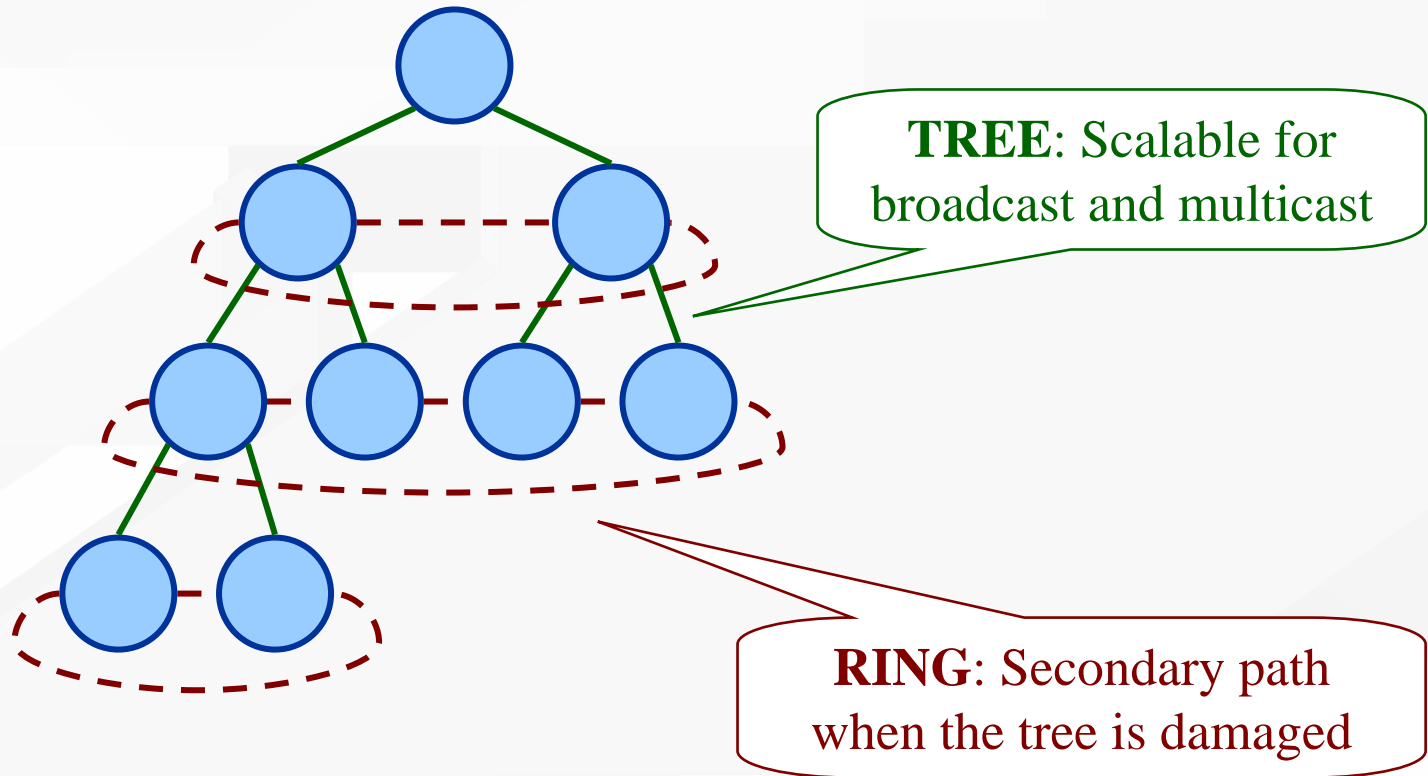
Design

Verification

Results

Conclusion

- » Based on k-ary sibling tree
 - » K is number of fan-out ($k \geq 2$)



Introduction

Background

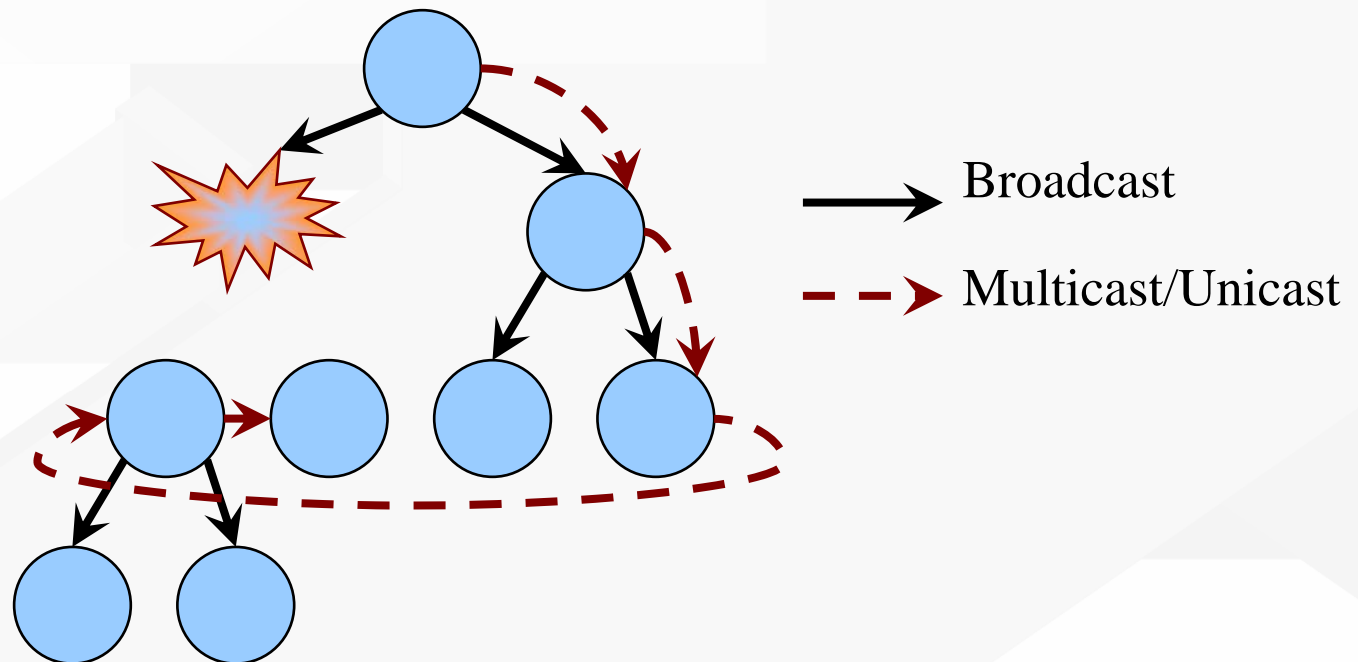
Design

Verification

Results

Conclusion

- » Example : survives a failure
 - » A broadcast message is encapsulated in a multicast message sent from parent to children of a dead node.



Introduction

Background

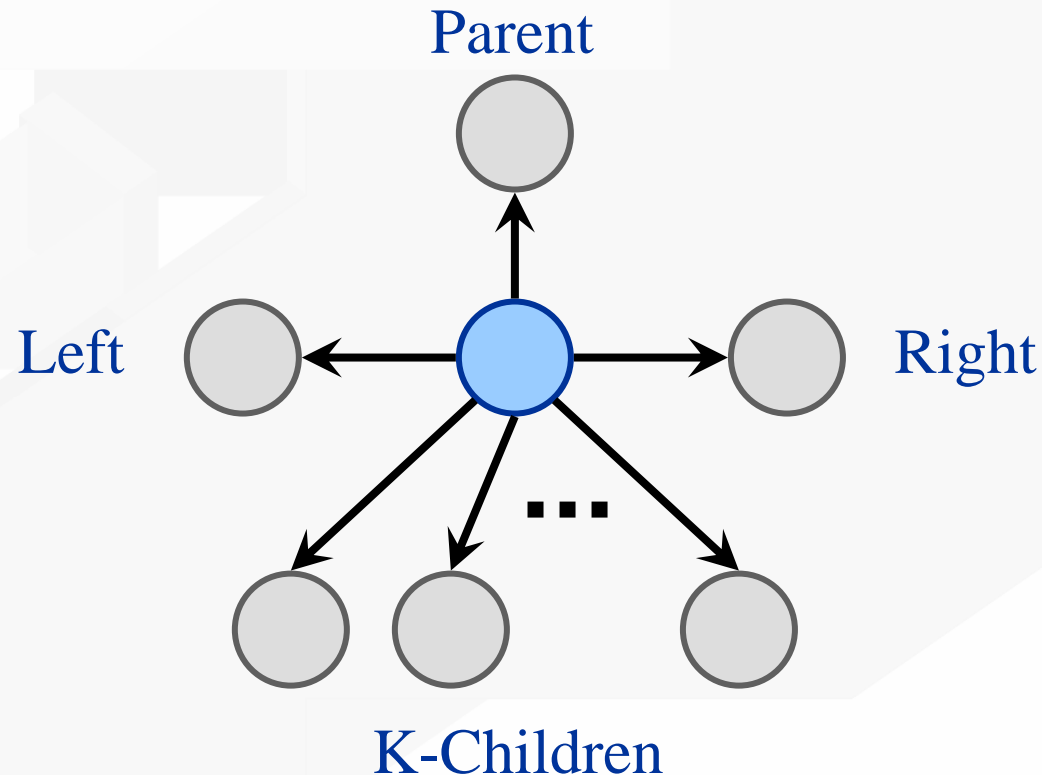
Design

Verification

Results

Conclusion

- » Low storage cost
 - » Each node needs to know
 - » the contact information of at most $k+3$ neighbors
 - » State of the link to its neighbors



Introduction

Background

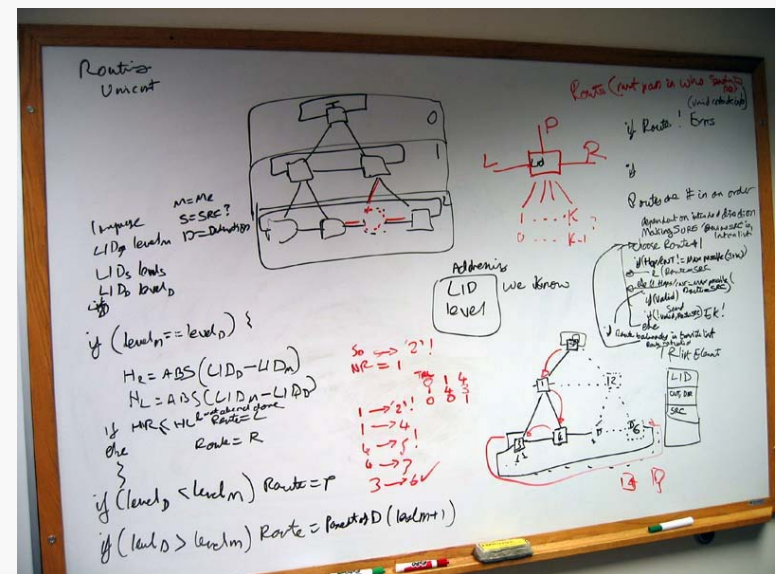
Design

Verification

Results

Conclusion

- » Protocol Specification
 - » Service Specification
 - » Environment Assumption
 - » Protocol Vocabulary
 - » Message Format
 - » Procedure Rules



Introduction

Background

Design

Verification

Results

Conclusion

- » Service Specification
 - » Deliver broadcast, multicast, unicast
 - » Normal circumstance
 - » Uses the k-ary tree to send messages
 - » Failure cases:
 - » Uses the neighbor to reroute messages
 - » Best effort routing

Introduction

Background

Design

Verification

Results

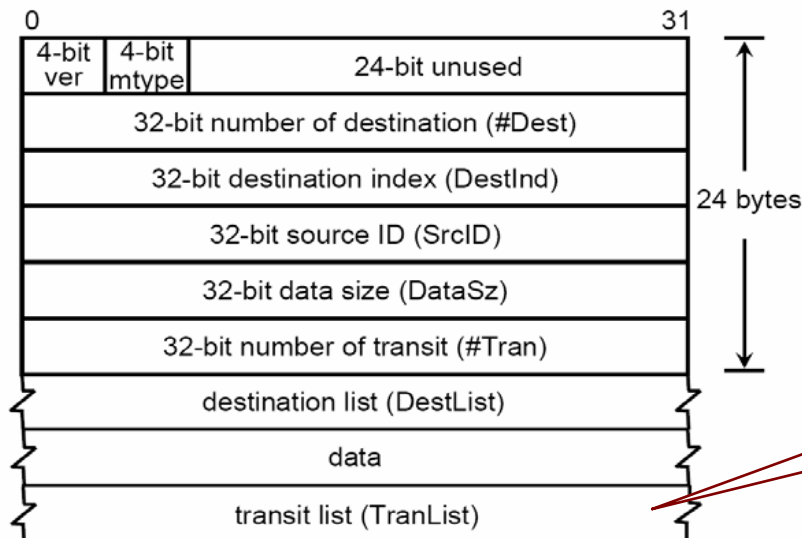
Conclusion

- » Environment Assumption
 - » Failures
 - » Assumes Fail-stop (rather than Byzantine)
 - » At least one neighbor of each node should be alive
 - » Unless allow each node to contact a directory service
 - » Transmission channel
 - » Can detect and recover from transmission error
 - » E.g. TCP, Reliable UDP
 - » Consequence: never lose a message
 - » Unless message is destroyed with a node before being pass on

» Protocol Vocabulary

- » Hello – Initialize messages (construct k-ary tree)
- » Mcast – Multicast messages (including Unicast)
- » Bcast – Broadcast messages

» Message Format



prevents message loop

Introduction

Background

Design

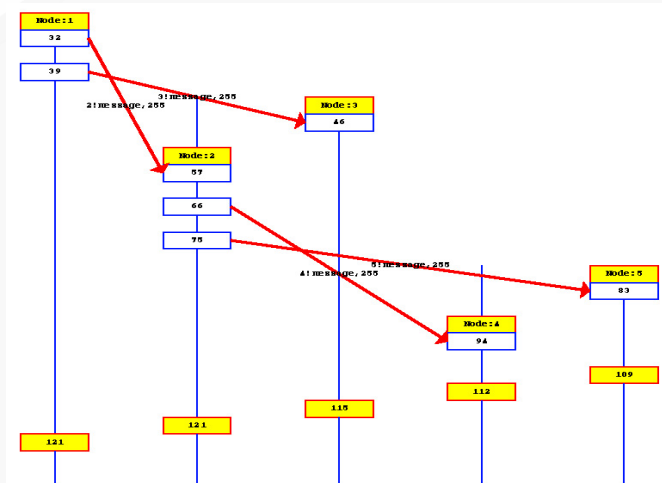
Verification

Results

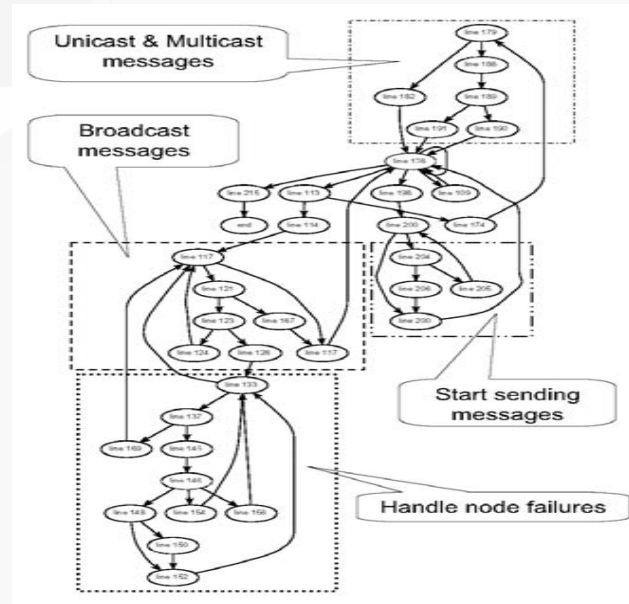
Conclusion

- » Procedure Rules:
 - » Initialization
 - » Register itself to the directory service
 - » Get its logical ID
 - » Send hello to Parent, Left
 - » (and to Right if the right most in each level)
 - » Routing (best effort)
 - » Bcast: send to all of its children
 - » If a child died: encapsulate in Mcast and reroute to its grand children
 - » Mcast: send to a valid neighbor (highest priority)
 - » Otherwise backtrack to sender
 - » ETC...

- » SPIN verification (and simulation) tool
 - » Model checker using automata-theoretical.
 - » Deadlocks, non-progress cycle, non-reachable state, etc.
 - » Provide counterexample in error cases.
 - » PROMELA (Process Meta Language)



- » Specifying the Protocol in PROMELA
 - » Model broadcast with exclusive channels
 - » Failures is simulated with non-deterministic selection ('if' selection construct)
 - » Speedup with 'atomic' construct



Introduction

Background

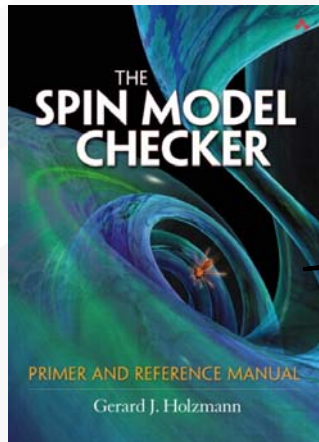
Design

Verification

Results

Conclusion

- » Verification Results
 - » No deadlock, livelock, invalid end state
 - » No unreachable codes and assertion violation



“I am SPIN and I approve this protocol”

Introduction

Background

Design

Verification

Results

Conclusion

» Routing Algorithms

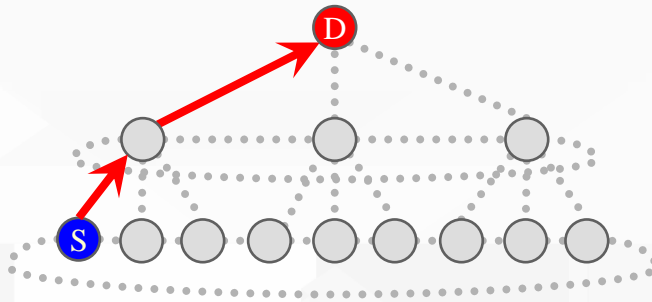
» 1) Basic

- » Fixed first hop based on static topology
- » Rule based method to estimate cost
- » Known locally failed links

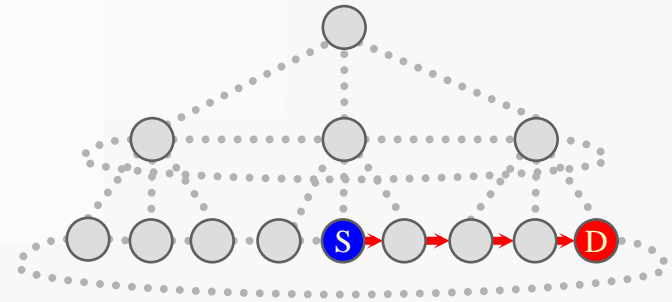
```
if my level = destination level then
    Send to left/right
else if my level > destination level then
    Send to my parent
else
    if my child is an ancestor of destination then
        Send to the child
    else
        Send to left/right who is closer to
        an ancestor of the destination in my level
```

Experimental Results

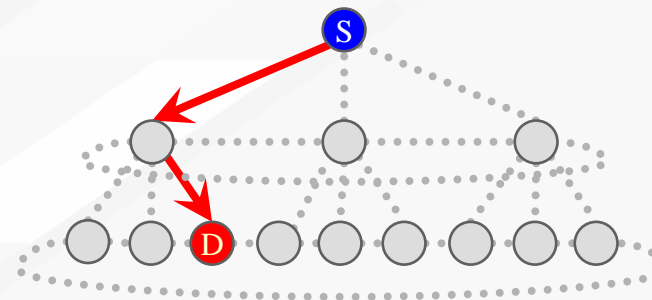
» Basic routing examples



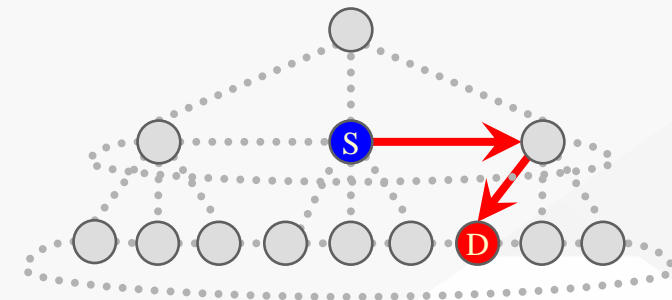
Destination Level above Sender Level



Destination Level = Sender Level



Case 1



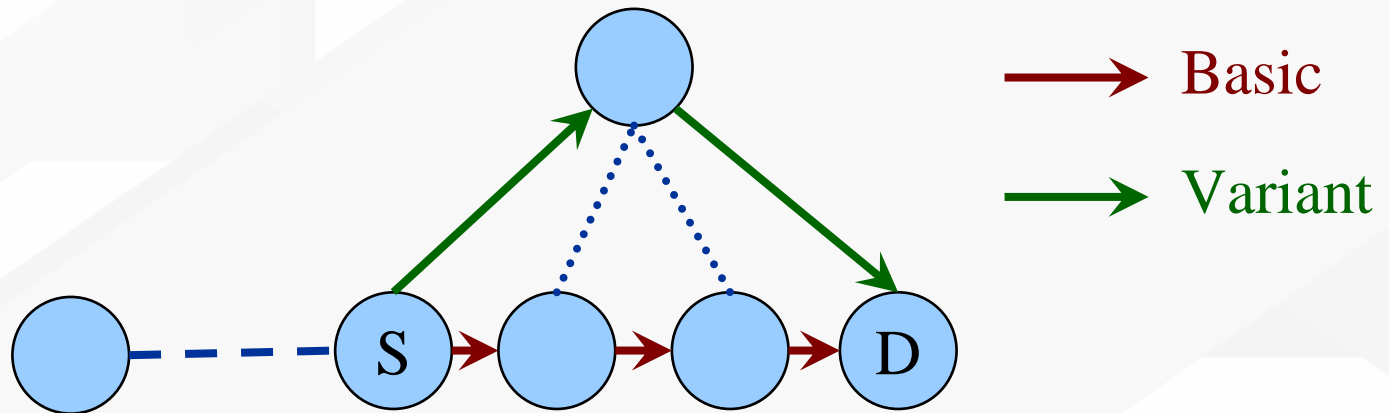
Case 2

Destination Level below Sender Level

» Routing Algorithms

» 2) Variant (of 1)

- » Based on ordering of current possible hops to shorten distance
- » (i.e. allows to go in a direction that does not toward destination)



Introduction

Background

Design

Verification

Results

Conclusion

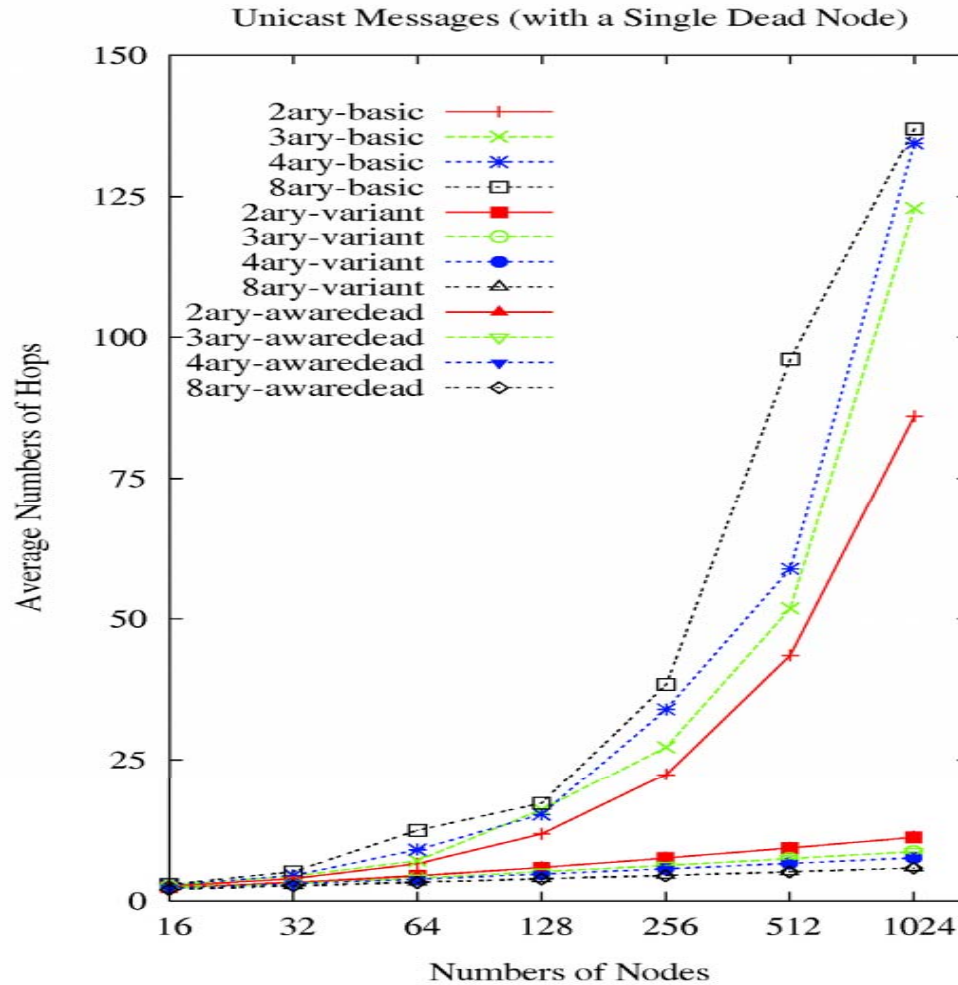
» Routing Algorithms

» 3) Breadth first search

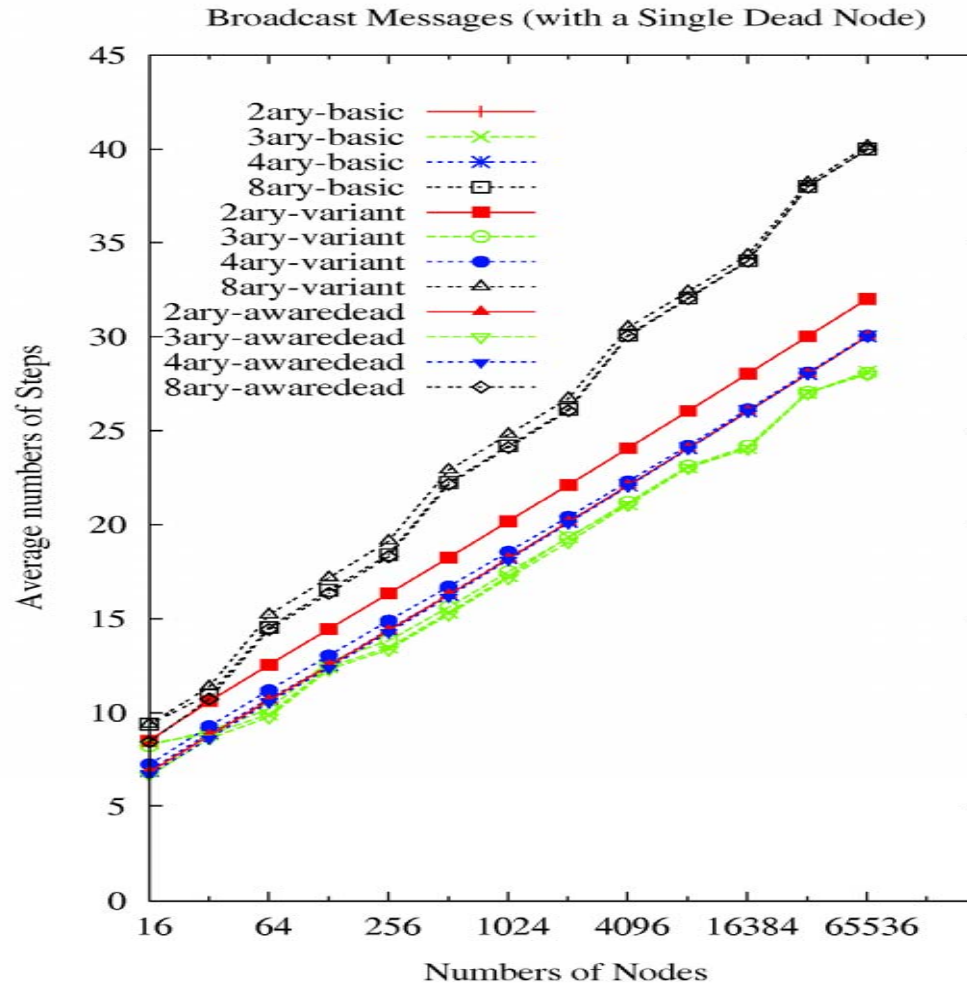
- » Graph-coloring which explore only alive nodes
- » Use knowledge of Previously detected dead nodes

» Note: more accurate, but time consumption





- Introduction
- Background
- Design
- Verification
- Results**
- Conclusion



Introduction

Background

Design

Verification

Results

Conclusion

- » Scalable and Fault-Tolerant Protocol
 - » Designed for parallel runtime environments
 - » Formally proven to work (normal and failure)

- » Future Work
 - » Protocol aware underlying network topology
 - » Add a function cost on each path
 - » Faster and more accurate re-routing algorithm
 - » Basic message distribution of Harness/Open RTE ala. FT-MPI/Open MPI

Please don't forget
the excursion at 4 PM 😊



Thank You / Danke schön

**For more information:
angskun@cs.utk.edu**