

Modeling and Verification of MPI based distributed software (using ASMs)

Igor Grudenić, Nikola Bogunović

Department of Electronics, Microelectronics, Computer and Intelligent Systems
Faculty of Electrical Engineering and Computing, Zagreb, Croatia

<http://www.fer.hr>

Presentation overview

- **Motivation**
- **Formal verification – three interesting properties**
- **Introduction to ASMs (really short one)**
- **Using ASMs to model basic MPI communication**
- **Mutual exclusion case study**
- **Conclusion & Future work**

- **Complex distributed systems are hard for design and analysis**
- **Typical problems**

	Synchronous communication	Asynchronous communication
Parallelism (scalability)	HARD	EASY?
Deadlock	POSSIBLE	NOT POSSIBLE
System complexity	MODERATE	HUGE

Formal verification

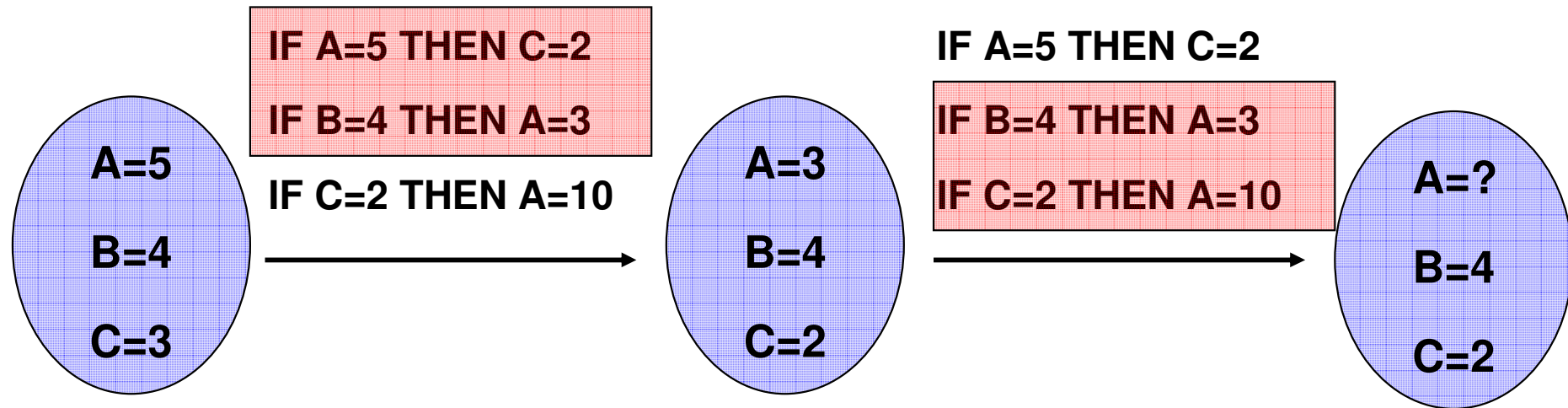
- **Checks whether system conforms to specification**
- **Theorem proving & model checking**
- **Model checking with SPIN (Siegel 2005)**
 - theorems and manipulation to avoid state explosion
 - three interesting properties for model checking and distributed systems
 - Safety (e.g. deadlock absence)
 - Liveness
 - Fairness

- **State machines – known concept**
- **Abstract state machines – similar to state machine but**
 - You need to define only the initial state and give the ‘recipe’ how to evolve (no enumeration)



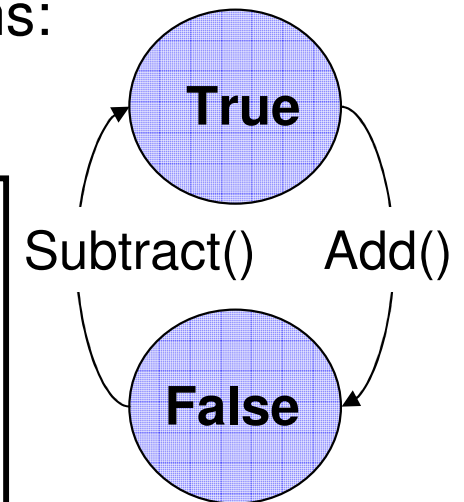
- Math:
 - Initial state (and every other state it evolved to) is a static algebra (set X , signature $\Sigma = \{f \mid X^r \rightarrow X\}$)
 - State evolution:
 - If **Condition** then **Update**
 - Update:** $f(t_1, \dots, t_n) = t, (f \in \Sigma \quad t_x, t \in X)$

Introduction to ASM (example)



- **ASML (abstract state machine language) – object oriented**
- **Model exploration – Spec Explorer**
- **How to define MPI processes ?**
 - every process is instance of the special *process class*
 - Every method in the process class is considered as atomic action
 - Example: process that executes sequence of actions:
Add(),Subtract(),Add(),...

```
class process()  
    var state as Boolean=true  
  
Add()                Subtract()  
    require state=true    require state=false  
    state:=false        state:=true
```

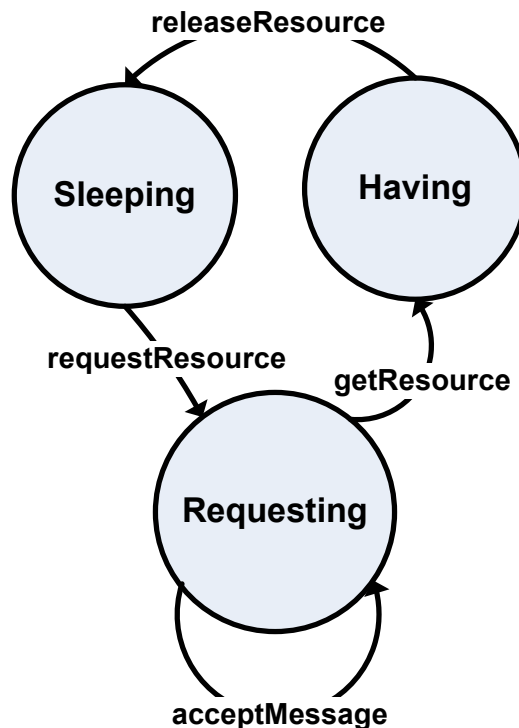


- Encoding MPI communication

	BLOCKING	NON BLOCKING
SEND	<pre> Blocking Send() require state=DESIRED STATE receiver.messageQueue+= [Message] state=MESSAGE SENT </pre>	<p>Make the same transfer as blocking but only when MPI_WAIT is posted?</p>
RECEIVE	<pre> AnnounceMessageReceive() require state=DESIRED STATE me.announces+={Announce()} state:=DESIRED STATE + RECEIVE ANNOUNCED ReceiveMessage() require messageQueue contains message (BL)require DESIRED STATE + RECEIVE ANNOUNCED applicationBuffer=getMessage(messageQueue) state:=MESSAGE RECEIVED </pre>	

Mutual exclusion – case study

- Lamport 1978, only non-blocking communication
- Three properties (**safety**, **liveness**, fairness)



requestResource ()

```

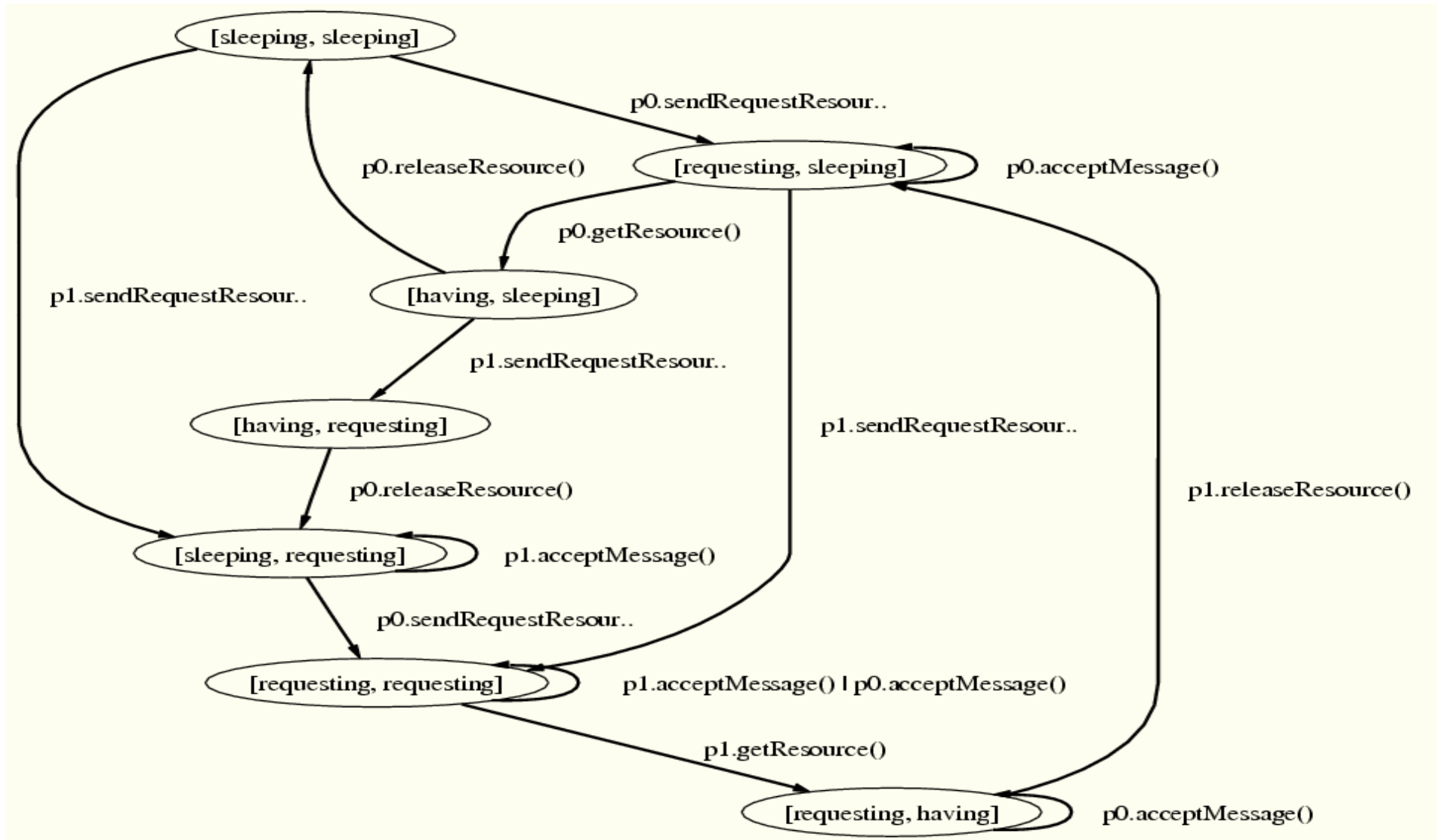
require state=tState.sleeping
state:=tState.requesting
foreach iProcess in processes
    where iProcess<>me
    iProcess.messageQueue+=[NewRequest]
    requestQueue+=[NewRequest]
    clock+=1, nOfReceivedAck:=0
  
```

acceptMessage ()

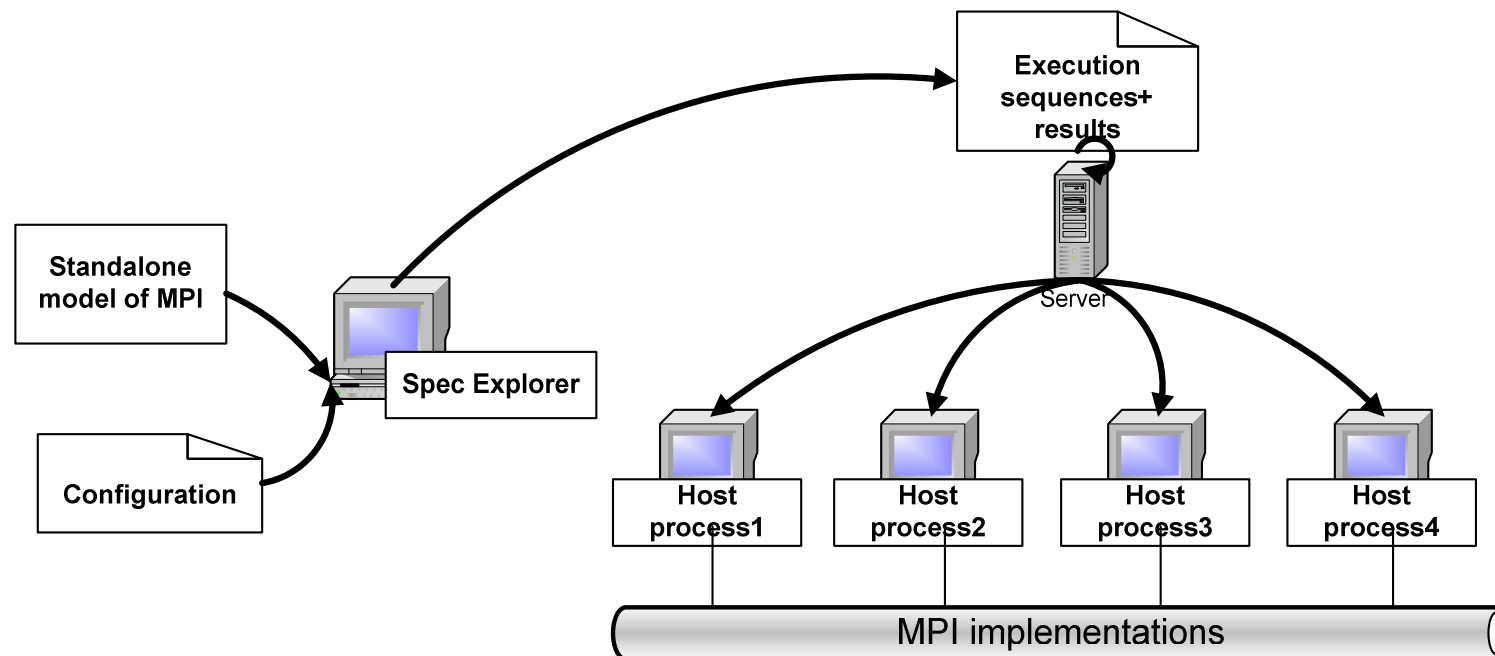
```

require Size(messageQueue)>0
require state=tState.requesting
messageQueue:=Tail(messageQueue)
clock:=max(messageQueue(0).clock+1, clock+1)
match messageQueue(0).messageType
  REQ: requestQueue+={messageQueue(0)}
      sender.messageQueue+=[NewAcknowledge]
  REL: requestQueue--=request(sender)
  ACK: nOfReceivedAck+=1
  
```

Mutual exclusion – exploration & verification



- **Standalone model of MPI (only basic communication)**
- **Configurations (made several by hand now)**
- **Host processes (support execution of MPI functions supported by the model)**



Conclusion

- **ASMs are powerful for capturing requirements**
- **ASMs can give an insight to a distributed system**
- **Some basic (but important) properties can be verified using available tools**
- **Questions?**